

MAT 312/AMS 351 Spring 2014 Review for Final (Mon-Wed-Fri section)

(use the reviews for Midterms 1 and 2 plus the following)

§5.4 and “Notes on binary codes.” Understand that a binary code of length n is a subset C of the abelian group $\mathbf{Z}_2 \times \dots \times \mathbf{Z}_2$ (n times), also written as \mathbf{B}^n (sequences of length n of 0s and 1s; addition is vector addition: componentwise, mod 2). The code C is *linear* if C is a subgroup of \mathbf{B}^n . (Since mod 2 every word is its own inverse, the condition amounts to C being closed under addition). Understand the (Hamming) distance p.234 between codewords and for a linear code know how to reckon the usefulness of a code (for error detection &/or correction), i.e. the minimum distance between codewords, by inspection of the set of code-words p.237. [After p.237 switch to material in the “Notes.”]

Interpreting codewords as vectors, understand how a homomorphism (a linear transformation) $h : \mathbf{B}^n \rightarrow \mathbf{B}^m$ defines a linear code C_h by $C_h = \{\mathbf{x} \in \mathbf{B}^n | h(\mathbf{x}) = \mathbf{0}\}$, the set of all n -tuples which h sends to the zero m -tuple in \mathbf{B}^m . (This is called the *kernel* of h). Understand how h can be represented by right-multiplication by a matrix H : the m -vector $h(\mathbf{x})$ is the matrix product $\mathbf{x}H$; the first row of H is the m -vector $h(1, 0, 0, \dots, 0)$, the second row is $h(0, 1, 0, \dots, 0)$, etc.); we can call the code C_H . Know how to deduce the error-detecting or error-correcting properties of C_H by inspection of the matrix H .

§6.1, 6.2 Understand the similarity between the divisibility of polynomials $s(x), t(x), \dots$ with coefficients in a field (the real numbers, \mathbf{Z}_2 , etc.) and the divisibility of integers. Be able to carry out the division algorithm (“long division”) for polynomials, giving a quotient and a remainder. Be able to carry out the Euclidean Algorithm to calculate a greatest common divisor $d(x)$ of $s(x), t(x)$ and to write $d(x)$ as a polynomial linear combination of $s(x)$ and $t(x)$. Note one difference: a polynomial $p(x)$ has a linear factor $(x - \alpha)$ if and only if $p(\alpha) = 0$. This is very useful in finite fields, since there are only finitely many possible α .

§6.3 Understand the definition of *irreducible* p.273; the distinction between irreducible and *prime* is not important in this context. Understand the proof of Theorem 6.3.4 (every polynomial can be written as a product of irreducibles) and the difference from Theorem 1.3.3 (unique factorization for integers): an irreducible factor is only determined *up to a nonzero multiplicative constant*. When the coefficient field is \mathbf{Z}_2 this difference does not manifest itself since the only nonzero constant is 1. Understand Examples 1 and 2 on p.277 completely.

§6.4 Understand that polynomial congruence classes are defined, and have many properties like, congruence classes of integers *mod* m . In particular, understand that when a polynomial $p(x)$ is irreducible, every nonzero congruence class *mod* $p(x)$ has a multiplicative inverse: Proposition 6.4.3, Example 2 p.281 and continuing in the Example on p.282. Be familiar with the examples worked out in class:

- in $\mathbf{Z}_2[X]$ using $p(x) = x^2 + x + 1$, $p(x) = x^3 + x + 1$
- in $\mathbf{Z}_3[X]$ using $p(x) = x^2 + 1$
- in $\mathbf{R}[X]$ using $p(x) = x^2 + 1$
- in $\mathbf{Q}[X]$ using $p(x) = x^2 - 2$

Be able to calculate products and inverses of equivalence classes in these and similar cases.